

# Complex-Network Modelling and Inference

## Lecture 7: Graph features

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# Section 1

## Graph features/metrics

# Graph Notation

- The network is defined by the *graph*,

$$G(N, E)$$

- We will assume (unless stated) that it is undirected.
- By default label the nodes  $\{1, 2, \dots, n\}$

# Graph Features/Metrics

- When graphs are small, we can draw them, and look at them, and it's still hard to assess them, *e.g.*, isomorphism
- When they are large it is impossible to visually assess them
- But we still need ways to analyse them
  - ▶ categorise
  - ▶ predict
  - ▶ assess for unusual characteristics
- There is a growing trend to do so using a common set of numbers called variously
  - ▶ statistics
  - ▶ metrics
  - ▶ features

# Graph Features/Metrics

There are two type of metrics/features

- Local (to the nodes)
  - ▶ node degree
  - ▶ local clustering coefficient
  - ▶ centrality (various versions)
  - ▶ eccentricity
- Local (to a pair of nodes)
  - ▶ (shortest path) distance
- Global (for the whole network)
  - ▶ average node degree and degree distribution
  - ▶ radius, average distance and diameter
  - ▶ global clustering coefficient
  - ▶ assortativity/homophily
  - ▶ graph spectrum

## Section 2

# Node degree distributions

# Neighbourhood and node degree

## Definition

The *neighbourhood* of node  $i$  is defined by

$$N_i = \{j \mid (i, j) \in E\},$$

*i.e.*, the set of nodes adjacent to  $i$ .

## Definition

The *node degree*  $k_i$  is the number of neighbours of node  $i$ , *i.e.*,

$$k_i = |N_i|.$$

## Definition (Alternative definition)

The *node degree*  $k_i$  is the number of edges incident to  $i$ , *i.e.*,

$$k_i = \left| \{(i, j) \mid (i, j) \in E\} \right|.$$

# Global node-degree statistics

- An often used statistic/feature of a graph is its *average node degree*

$$\bar{k} = \langle k \rangle = \frac{1}{|N|} \sum_{i \in N} k_i = \frac{2|E|}{|N|}.$$

the last result by the Handshake theorem.

- More generally the *node degree distribution*  $p_k$  gives the probability that a node has degree  $k$  (or relative frequency)
  - ▶ an  $r$ -regular graph has

$$p_k = \begin{cases} 1, & \text{for } k = r, \\ 0, & \text{otherwise.} \end{cases}$$



# Friendship paradox [Fel91]

*Friendship paradox = your friends have more friends than you*

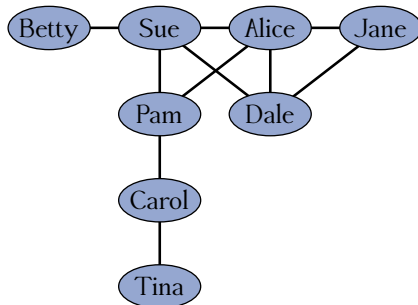
- Actually the theorem is statement about averages  
“On average, your friends will have more friends than you.”
- Stated mathematically

$$E[k_i] \leq E[k_{neighbours(i)}].$$

- Some versions are about “most” people, *i.e.*, most people’s friends have more friends than them
- Intuition is that sampling the node-degree distribution by looking at friends artificially biases the high-degree nodes because they are friends more often.

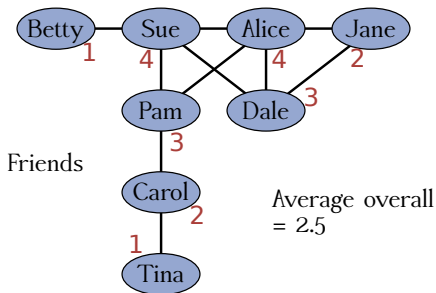
# Friendship paradox example [Fel91]

Marketville High School Girls (subgraph)



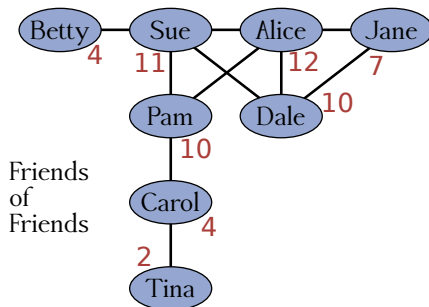
# Friendship paradox example [Fel91]

Marketville High School Girls (subgraph)



# Friendship paradox example [Fe191]

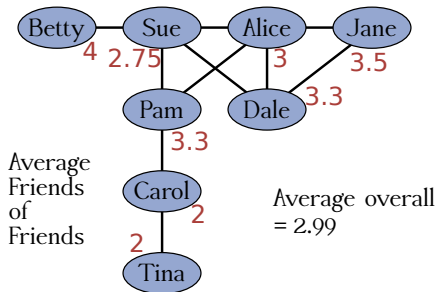
Marketville High School Girls (subgraph)



Friends  
of  
Friends

# Friendship paradox example [Fel91]

Marketville High School Girls (subgraph)



# Friendship paradox

## Proof.

The average number of friends is

$$\bar{k} = \frac{1}{n} \sum_{i=1}^n k_i.$$

When we calculate the number of friends of friends  $k_i^{(2)}$ , Feld [Fel91] argued that each individual is “a friend  $k_i$  times and has  $k_i$  friends, so that individual contributes ... a total of  $k_i^2$  friends’ friends”. Thus the total friends’ friends is

$$\sum_{i=1}^n k_i^{(2)} = \sum_{i=1}^n k_i^2,$$

and we average this over the total number of friends, *i.e.*,  $\sum_i k_i$ , to get

$$\overline{k^{(2)}} = \frac{\sum_{i=1}^n k_i^2}{\sum_{i=1}^n k_i}.$$

# Friendship paradox

Proof.

Standard result

$$\text{Var}(X) = \mathbb{E}[X^2] - \mathbb{E}[X]^2.$$

Re-arranging we get

$$\frac{\mathbb{E}[X^2]}{\mathbb{E}[X]} = \mathbb{E}[X] + \frac{\text{Var}(X)}{\mathbb{E}[X]}.$$

In our context

$$\overline{k^{(2)}} = \frac{\sum_{i=1}^n k_i^2}{\sum_{i=1}^n k_i} = \mathbb{E}[k_i] + \frac{\text{Var}(k_i)}{\mathbb{E}[k_i]},$$

and we know the mean and variance are  $\geq 0$  so

$$\overline{k^{(2)}} \geq \bar{k}.$$



# Friendship paradox

## Proof (part 2)

- I find the first part a little hand-wavy
- There is a nice little lesson to learn in doing it mathematically

### Proof.

We can write the number of friends of  $i$  using the adjacency matrix  $A = [a_{ij}]$

$$k_i = \sum_j a_{ij} = \sum_j a_{ji}.$$

We can similarly write the number of friends of friends for  $i$  using the adjacency matrix by considering that a friend  $j$  of  $i$  will be reached by  $a_{ij}$ , so

$$k_i^{(2)} = \sum_j a_{ij} k_j.$$





## An aside

The number  $i$ 's friends' friends can be seen as

$$k_i^{(2)} = \sum_j a_{ij} k_j = \sum_j a_{ij} \sum_k a_{jk} = \sum_k \sum_j a_{ij} a_{jk}$$

Which is just the matrix squared, *i.e.*, if  $A^2 = [a_{ij}^{(2)}]$  then we can think of this as

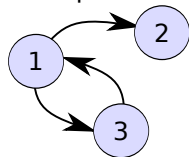
$$a_{ik}^{(2)} = \sum_j a_{ij} a_{jk} = \text{the number of 2 hop paths from } i \text{ to } j,$$

and  $k_i^{(2)}$  is the sum over the possible end-points of such paths.

*The main point is  $A^2$  contains the number of two-hop paths between each pair of nodes.*

## An aside: example

Example directed graph



$$A = \begin{pmatrix} 0 & 1 & 1 \\ 0 & 0 & 0 \\ 1 & 0 & 0 \end{pmatrix}$$

Now

$$A^2 = \begin{pmatrix} 1 & 0 & 0 \\ 0 & 0 & 0 \\ 0 & 1 & 1 \end{pmatrix}$$

So there are

- exactly one 2-hop path from 1 to 1: 1-3-1
- exactly one 2-hop path from 3 to 2: 3-1-2
- exactly one 2-hop path from 3 to 3: 3-1-3
- no other 2-hop paths

# Friendship paradox

Proof.

$$k_i^{(2)} = \sum_j a_{ij} k_j.$$

So the total number of friends' friends is

$$\sum_i k_i^{(2)} = \sum_i \sum_j a_{ij} k_j.$$

Principle of perversity (of sums) leads us to change the order of summation

$$\sum_i \sum_j a_{ij} k_j = \sum_j \sum_i a_{ij} k_j = \sum_j k_j \sum_i a_{ij} = \sum_j k_j^2.$$

Hence

$$\sum_i k_i^{(2)} = \sum_j k_j^2.$$



## Section 3

# Homophily and Assortativity

# Homophily

*Birds of a feather, flock together.*

- Homophily expresses the idea that many relationships (that we might express in a graph) are more likely between similar entities.
- Many studies have confirmed it in many contexts
  - ▶ characteristics: age, gender, class, geography, ...
  - ▶ relationships: collaboration, friendship, ...
- In random graphs (see next week) assume the probability of a link depends on similarity of characteristics of the nodes
  - ▶ we'll define in terms of  $e_{ij}$  = fraction of edges that connect a vertex of type  $i$  to one of type  $j$ .

# Homophily and assortativity

- Assortative mixing, or just *assortativity* expresses homophily between nodes based on their node degree.
- The definition is slightly circular: edges are more common between nodes with similar numbers of edges ...
  - ▶ But we can work with that
- Measure using Pearson correlation coefficient of “remaining” degrees at either ends of a random edge.
  - ▶ start with the idea of a correlation of nodes on a random link
  - ▶ but eliminate the link in question  $\Rightarrow$  remaining
  - ▶  $e_{jk}$  is the joint probability distribution of the remaining degree at either end of a randomly chosen link

## Remaining degree distribution

- The degree distribution of a node reached by a random link

$$q'_k = \frac{k p_k}{\sum_j j p_j}, \quad k = 1, 2, \dots$$

- The *remaining degree distribution*, ignores the link we came in on

$$q_k = q'_{k+1} = \frac{(k+1)p_{k+1}}{\sum_j j p_j}, \quad k = 0, 1, 2, \dots$$

- $\sigma_q^2$  variance of of distribution  $q_k$

$$\sigma_q^2 = \sum_k k^2 q_k - \left[ \sum_k k q_k \right]^2$$

- $q_k$  is the marginal distribution of  $e_{jk}$ , i.e.,

$$q_k = \sum_j e_{jk}$$

## Metric 2: assortativity

- Assortativity

$$r = \frac{\sum_{j,k} jk(e_{jk} - q_j q_k)}{\sigma_q^2}$$

- $r$  is the Pearson correlation coefficient of remaining degrees at either ends of a random edge.

$$-1 \leq r \leq 1$$

- cases

- ▶  $r$  near 1 means high degree nodes often connect to high degree nodes
- ▶  $r$  near -1 means high degree nodes often connect to low degree nodes



# Evaluating assortativity

In a real network, we evaluate  $r$  by taking

$$\hat{r} = \frac{w \sum_{e \in E} j_e k_e - [w \sum_{e \in E} (j_e + k_e)/2]^2}{w \sum_{e \in E} (j_e^2 + k_e^2)/2 - [w \sum_{e \in E} (j_e + k_e)/2]^2},$$

where

$$w = |E|^{-1}$$

$j_e$  = degrees of vertex at one end of the edge  $e$

$k_e$  = degrees of vertex at other end of the edge  $e$

# Further reading I



Scott L. Feld, *Why your friends have more friends than you do*, *American Journal of Sociology* **96** (1991), no. 6, 1464–1477.