

# Information Theory and Networks

## Lecture 16: Gambling and Information Theory

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# Part I

## Gambling and Information Theory

If fighting is sure to result in victory, then you must fight,  
even though the ruler forbid it;  
If fighting will not result in victory, then you must not fight  
even at the ruler's bidding.

*Sun Tzu, The Art of War, Chapter 10, 23*

# Section 1

## Horse Racing

# Fixed-Odds Horse Racing

- Pool of money betting on horses
  - ▶ odds: expressed as  $o$ -for-1 or  $(o - 1)$ -to-1
  - ▶ probability of success by probability of failure
  - ▶ assume no track take, no commissions
- What's the best strategy?
  - ▶ one-off bet
  - ▶ multiple ongoing bets, or *parlayed bets*

## Example

- Here, only bet on horse win (not other bets like place etc.)
- Odds are fixed by a bookie
- We use *o*-for-1 convention

Horse	Odds
1	10
2	2
3	20
4	5

# Betting Strategies

- One-off bet: all in
  - ▶ equivalent: maximizing arithmetic mean
- Parlayed bets: Kelly criterion
  - ▶ equivalent: maximizing geometric mean
- What happens with all-in for parlayed bets?
- Note: payout asymmetry most important
- **Make sure your capital survives before it can compound**

## Section 2

# The Kelly Criterion



# Some History

- Developed by J. L. Kelly at Bell Labs; Shannon reviewed
  - ▶ Texan tough guy, gunslinger, daredevil pilot and mathematician!
- Wirelines were used to transmit information between bookies
  - ▶ application: placing bets on horses

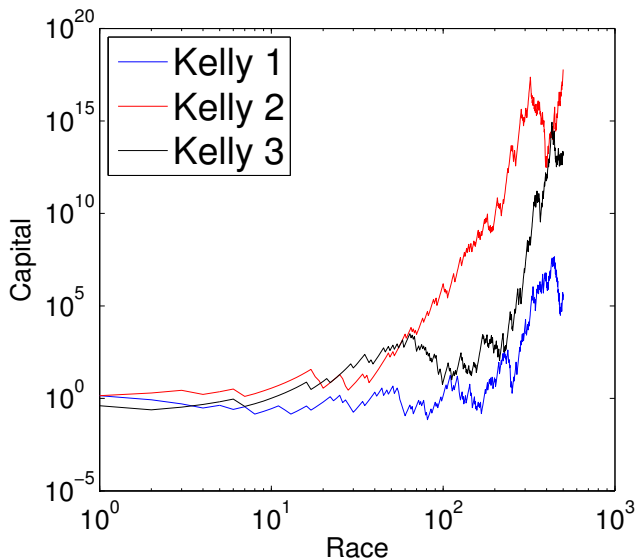
# Formulation

- Assume  $m$  horses, each with i.i.d. probability of winning  $p_i$
- Assume starting capital  $S_0 = 1$
- Odds:  $o_i$ , alternative  $(1 + r_i)$ ,  $r_i$  the rate of return
- Play for  $T$  races
  - ▶ allocate  $b_i$  fraction of capital on horse  $i$
  - ▶ capital at  $T$ :  $S_T = \prod_{t=1}^T \prod_{i=1}^m b_i o_i$
- Objective: assuming fully invested, choose allocation  $b_i \geq 0$ ,  $\sum_i b_i = 1$  to maximize  $S_T$

# Maximising Wealth Growth

- Assume  $T \rightarrow \infty$ 
  - ▶ maximise  $E[\sum_{i=1}^m \log b_i o_i]$  subject to constraints
  - ▶ doubling rate:  $W(\mathbf{b}, \mathbf{p}) := \sum_{i=1}^m p_i \log b_i o_i$
- Solution: the Kelly criterion, or log-optimal wealth growth
  - ▶ answer:  $b_i^* = p_i$ , proportional gambling (for fair odds)
  - ▶ solve using standard KKT conditions, or log-sum inequality
- Nature of solution will depend on odds: see [CT91, Exercise 6.2]

# Example Run of Kelly's Strategy

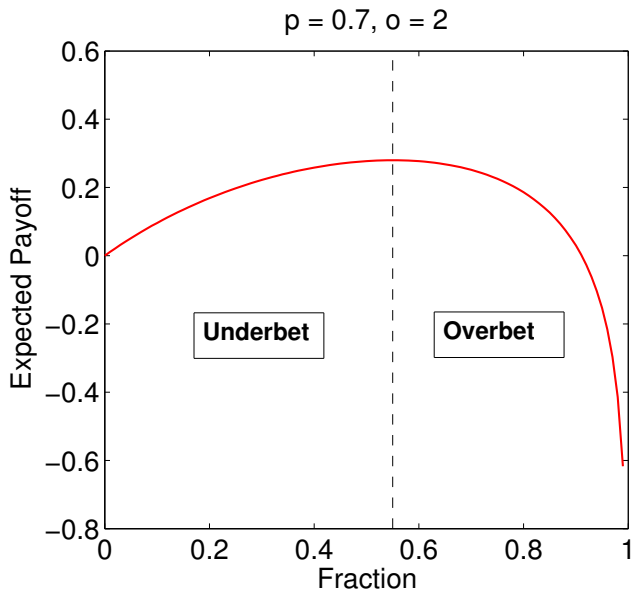


# A Simple Bet

- Say a biased coin toss, win if heads, lose if tails
  - ▶ heads with probability  $p$ ,  $q$  otherwise
  - ▶ each round, add \$1 to bet
- Odds:  $o$ -for-1 (remember: win-lose event)
- Kelly solution:  $b^* = \frac{op-q}{o} = \frac{p(o+1)-1}{o}$ 
  - ▶ what does it mean if  $o = q/p$ ?
  - ▶ what does it mean when  $b^* < 0$  ( $o < q/p$ )?
  - ▶ what about  $b^* > 1$ ?
- A simple way to remember (for two events)

$$b^* = \frac{\text{edge}}{\text{odds}}$$

# Simple Bet: Payoff



# Simple Bet: Under and Overbetting

- There is no gain in overbetting: growth decreases, risk increases
- Sweet spot: full Kelly for maximum wealth growth
- In practice, partial Kelly more applicable, i.e.  $\alpha b_i^*$ 
  - ▶ with  $\alpha$  fraction, only  $\alpha^2$  volatility
  - ▶ more robust to error in estimating returns
  - ▶ lower wealth growth compared to full Kelly

## Section 3

# Downsides



# Caveats

- Strategy is guaranteed to beat any other strategy on wealth growth
- BUT Strategy is asymptotically optimal: assume playing forever
- No guarantee to win in the short term (or at all), just the best chance
- Psychologically unsettling: imagine capital dropping 60% right before tripling!
  - ▶ partial Kelly strategies trade smoothness with growth rate
- Guaranteed not to go to ruin
  - ▶ BUT assumes capital infinitely divisible
  - ▶ capital could be  $10^{-10}$  but hey, at least not bankrupt!
  - ▶ can show  $\lim_{T \rightarrow \infty} P(S_T > \epsilon) = 0$ , for any  $\epsilon > 0$
- Assumes know the probability of winning: not true in real life
  - ▶ again, half Kelly strategies help: gives a safety margin
  - ▶ estimation methods (e.g. maximum entropy, shrinkage)

# Criticism from Modern Finance

- Kelly criterion assumes maximizing growth rate exponent
- Called the log-utility function in finance
- **Criticism 1:** not everybody would want to maximise growth rate exponent
  - ▶ does not take into account risk-averseness (or “sleep test”)
  - ▶ definition of risk in finance: volatility
  - ▶ different utilities for different folks
- **Criticism 2:** time horizon, as discussed, need very long term
- Counter-argument: not many people want to do with less money
- “Money can’t buy you happiness, but love can’t get you a Ferrari.”

# Approximation of the Stock Market

- Suppose  $m$  risky assets, each with random “odds”  $r_i$  in one investment period
- One asset with return  $r_0$  is deterministic
- Assume starting capital  $S_0 = 1$
- The return vector  $\mathbf{r}$ , with  $\boldsymbol{\mu}_r = E[\mathbf{r}]$ ,  $\Sigma = E[(\mathbf{r} - r_0\mathbf{1})(\mathbf{r} - r_0\mathbf{1})^T]$ 
  - ▶  $\Sigma$  is full rank
  - ▶ correlations apply only “spatially”
- Derive the optimal allocation  $\mathbf{b}$  to optimise the wealth doubling rate
  - ▶ optimise  $E[\log(r_0 + \mathbf{b}^T(\mathbf{r} - r_0\mathbf{1}))]$
- Assume no constraints on  $\mathbf{b}$
- For what return distribution is this allocation optimal?

# Further reading I



Thomas M. Cover and Joy A. Thomas, *Elements of information theory*, John Wiley and Sons, 1991.